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**END OF SEMESTER ASSESSMENT PAPER**

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MODULE TITLE: Optimisation

DURATION OF EXAMINATION:  $2\frac{1}{2}$  hours

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GRADING SCHEME: 20% +80%

EXTERNAL EXAMINER: Prof. J. King

**INSTRUCTIONS TO CANDIDATES: Answer four questions correctly for full marks; 80%. See the Appendix at the end of the paper for some useful results.**

- 1 (a) Suppose that a function  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  is  $C^1$  on  $\mathbb{R}^n$  and has gradient  $\nabla f(\mathbf{x}) \equiv \mathbf{g}(\mathbf{x})$ . If a vector  $\mathbf{p}$  satisfies the condition  $\mathbf{p}^T \mathbf{g}(\mathbf{x}) < 0$ ,  $\mathbf{p}$  is called a “descent direction” for  $f$  at  $\mathbf{x}$ . Explain carefully the importance of this definition for optimisation algorithms, justifying any statement that you make. 4
- (b) Suppose that  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  is  $C^1$  and  $\mathbf{p}_k$  is a descent direction at  $\mathbf{x}_k$ . Assume that  $f$  is bounded below along the ray  $\{\mathbf{x}_k + \alpha \mathbf{p}_k | \alpha > 0\}$ . Show that if  $0 < c_1 < c_2 < 1$  there exist step lengths  $\alpha$  satisfying the Wolfe conditions; Eqs. 10a and 10b in Appendix A. 12
- (c) Draw a sketch to illustrate your proof. 4
- (d) Show that if  $c_1 > \frac{1}{2}$ , a line search satisfying the Wolfe conditions would exclude the minimiser of a strictly convex quadratic objective function so the restriction  $c_1 \leq \frac{1}{2}$  is reasonable. 5
- Hint: write  $\phi(\alpha) = \phi_0 + \alpha \phi'_0 + \frac{1}{2} \alpha^2 \phi''_0$ , find the  $\alpha$ -range determined by the first Wolfe condition and compare with  $\alpha^*$ , the  $\alpha$ -value that minimises  $\phi(\alpha)$ . Take  $\phi''_0 > 0$  as given  $f$  strictly convex.
- 2 (a) Prove Zoutendijk’s Theorem (Theorem 1 in Appendix C ) 15
- (b) Explain briefly the significance of the result. 1
- (c) Suppose that search directions  $\mathbf{p}_k$  are generated using a Newton-like method:  $\mathbf{p}_k = -\mathbf{B}_k^{-1} \mathbf{g}(\mathbf{x}_k)$  where  $\mathbf{B}_k$  is symmetric and positive definite. For any matrix  $\mathbf{A}$  let  $\|\mathbf{A}\|$  be the matrix 2-norm, equal for real symmetric matrices to the absolute value of the largest eigenvalue of  $\mathbf{A}$ . Show that if  $\|\mathbf{B}_k\| \|\mathbf{B}_k^{-1}\| \leq M$  for all  $k$  then  $\cos \theta_k \geq 1/M$  where  $\theta_k$  is defined in Appendix C. 8
- (d) Use Zoutendijk’s Theorem above to show that in this case  $\lim_{k \rightarrow \infty} \|\mathbf{g}(\mathbf{x}_k)\| = 0$ . 1

3 The Fletcher-Reeves (FR) version of the non-linear conjugate gradient algorithm (Alg. 1) is given in Appendix E.

- (a) Suppose that the algorithm is implemented with a step length  $\alpha_k$  that satisfies the strong Wolfe conditions with  $0 < c_2 < \frac{1}{2}$  and that the norm of the gradient is bounded above. Assume Zoutendijk's Theorem (Theorem 1 in Appendix C) and the result that

$$-\frac{1}{1-c_2} \leq \frac{\mathbf{p}_k^\top \mathbf{g}_k}{\|\mathbf{g}_k\|^2} \leq \frac{2c_2-1}{1-c_2}, \text{ for all } k = 0, 1, \dots \quad (1)$$

and prove that the FR cgm has the global convergence property — i.e. that

$$\liminf_{k \rightarrow \infty} \|\mathbf{g}_k\| = 0.$$

Hints: (each step uses some or all of those preceding it). Use a proof by contradiction — so assume that  $\|\mathbf{g}_k\|$  is bounded below, say by  $\gamma > 0$ .

- (i) First use Eq. 1 above to show that

$$\frac{1-2c_2}{1-c_2} \frac{\|\mathbf{g}_k\|}{\|\mathbf{p}_k\|} \leq \cos \theta_k \leq \frac{1}{1-c_2} \frac{\|\mathbf{g}_k\|}{\|\mathbf{p}_k\|}.$$

- (ii) Then use Zoutendijk's Theorem to show that  $\sum_{k=0}^{\infty} \frac{\|\mathbf{g}_k\|^4}{\|\mathbf{p}_k\|^4}$  converges.

- (iii) Use the strong version of the second Wolfe condition (Eq. 11 in App. A) and the bound on  $\cos \theta_k$  above to show that  $|\mathbf{g}_k^\top \mathbf{p}_{k-1}| \leq \frac{c_2}{1-c_2} \|\mathbf{g}_{k-1}\|^2$ .

- (iv) Use the update rule for  $\mathbf{p}_{k+1}$  in the FR CGM algorithm (Appendix E) (replacing  $k+1$  by  $k$ ) to show that

$$\|\mathbf{p}_k\|^2 \leq c_3 \|\mathbf{g}_k\|^2 + \beta_k^2 \|\mathbf{p}_{k-1}\|^2$$

$$\text{with } c_3 = \frac{1+c_2}{1-c_2}.$$

- (v) Finally iterate this inequality to obtain a contradiction with the result of hint (iii) — you can assume that  $\|\mathbf{g}_k\|$  is bounded above, say by  $\bar{\gamma}$ .

- (b) Explain the significance of the result.

4 (a) Prove the Sherman-Morrison-Woodbury formula for the inverse of a matrix  $\mathbf{A} + \Delta\mathbf{A}$  — stated in Appendix F.

(b) Given the DFP update formula:

$$\mathbf{DFP} : \mathbf{H}_{k+1} = (\mathbf{I} - \gamma_k \mathbf{y}_k \mathbf{s}_k^T) \mathbf{H}_k (\mathbf{I} - \gamma_k \mathbf{s}_k \mathbf{y}_k^T) + \gamma_k \mathbf{y}_k \mathbf{y}_k^T, \quad (2)$$

where

$$\gamma_k = \frac{1}{\mathbf{y}_k^T \mathbf{s}_k}, \mathbf{s}_k = \mathbf{x}_{k+1} - \mathbf{x}_k \quad \text{and} \quad \mathbf{y}_k = \nabla f(\mathbf{x}_{k+1}) - \nabla f(\mathbf{x}_k),$$

show that the DFP-update formula may be written as  $\mathbf{H}_{k+1} = \mathbf{H}_k + \Delta \mathbf{H}_k$ , with

$$\Delta \mathbf{H}_k = \mathbf{R} \mathbf{S} \mathbf{R}^T,$$

where

$$\mathbf{R} = [\mathbf{y}_k \quad \mathbf{H} \mathbf{s}_k], \quad \mathbf{S} = \gamma_k \begin{bmatrix} 1 + \gamma_k \mathbf{s}_k^T \mathbf{H}_k \mathbf{s}_k & -1 \\ -1 & 0 \end{bmatrix}.$$

(c) Apply the SMW formula from part (a) to the formula  $\mathbf{H}_{k+1} = \mathbf{H}_k + \Delta \mathbf{H}_k$  from part (b) to derive the following equation for the update of the inverse Hessian approximation,  $\mathbf{J}_k$ , that corresponds to the DFP update of  $\mathbf{H}_k$  in Eq. 2;

$$\mathbf{DFP} - \mathbf{Inverse} \quad \mathbf{J}_{k+1} = \mathbf{J}_k - \frac{\mathbf{J}_k \mathbf{y}_k \mathbf{y}_k^T \mathbf{J}_k}{\mathbf{y}_k^T \mathbf{J}_k \mathbf{y}_k} + \frac{\mathbf{s}_k \mathbf{s}_k^T}{\mathbf{y}_k^T \mathbf{s}_k}. \quad (3)$$

Remember that for any  $2 \times 2$  matrix

$$\begin{bmatrix} \mathbf{a} & \mathbf{b} \\ \mathbf{c} & \mathbf{d} \end{bmatrix}^{-1} = \frac{1}{(\mathbf{a}\mathbf{d} - \mathbf{b}\mathbf{c})} \begin{bmatrix} \mathbf{d} & -\mathbf{b} \\ -\mathbf{c} & \mathbf{a} \end{bmatrix}.$$

5 Consider the general equality-constrained problem (18) as defined in Appendix H.

(a) Show that the penalty function  $F^k(\mathbf{x})$  defined in (17) in Appendix G has the property that the (unconstrained) minima  $\mathbf{x}_k$  of  $F^k$  converge to a local minimum  $\mathbf{x}^*$ .

(b) Prove the KKT first-order necessary conditions (20a) and (20b) in Appendix J for an **equality**-constrained problem .

(c) Finally, prove that when inequality constraints are introduced, the extra condition  $\lambda_i^* \geq 0$  for all  $i \in \mathcal{I}$  is necessary for  $\mathbf{x}^*$  to be an optimal solution of (19) in Appendix I. Define the functions  $\mathbf{c}_i^-(\mathbf{x}) = \min\{0, \mathbf{c}_i(\mathbf{x})\}$ ,  $j \in \mathcal{I}$ . Obviously  $\mathbf{c}_i^-(\mathbf{x}) \leq 0$ ,  $j \in \mathcal{I}$ .

Adapt (without repeating the details) your argument for part (b) to show that  $\lambda_i^* = -\lim_{k \rightarrow \infty} k \mathbf{c}_i^-(\mathbf{x}_k)$ ,  $i \in \mathcal{I}$ .

6 (a) Consider the equality-constrained Quadratic Program (QP):

$$\min_{\mathbf{x}} q(\mathbf{x}) = \frac{1}{2}\mathbf{x}^T\mathbf{Q}\mathbf{x} + \mathbf{x}^T\mathbf{d}, \quad (4a)$$

$$\text{subject to } \mathbf{a}_i^T\mathbf{x} = \mathbf{b}_i, \quad i = 1, \dots, k \quad (4b)$$

Let  $\mathbf{A}$  be the matrix s.t. the vectors  $\{\mathbf{a}_i\}_{i \in \mathcal{E}}$  are the columns of  $\mathbf{A}^T$ .

(i) Writing the set of  $k$  equality constraints (4b) as the matrix equation  $\mathbf{A}\mathbf{x} - \mathbf{b} = \mathbf{0}$ , show that if  $\mathbf{x}^*$  is a local minimum then the KKT conditions (Appendix J) require that there must be a vector  $\boldsymbol{\lambda}^*$  of Lagrange multipliers such that

2

$$\mathbf{Q}\mathbf{x} + \mathbf{d} - \sum_{i \in \mathcal{E}} \lambda_i^* \mathbf{a}_i = \mathbf{0}. \quad (5)$$

(ii) Writing  $\mathbf{x}^* = \mathbf{x}_0 + \mathbf{p}$ , where  $\mathbf{x}_0$  is any estimate of the solution and  $\mathbf{p}$  the required step to the solution, show that (6) and the constraints  $\mathbf{A}\mathbf{x} - \mathbf{b} = \mathbf{0}$  can be re-written as:

3

$$\mathbf{A}\mathbf{Q}^{-1}\mathbf{A}^T\boldsymbol{\lambda} = \mathbf{A}\mathbf{Q}^{-1}\mathbf{g} - \mathbf{r}, \quad (6)$$

$$\mathbf{Q}\mathbf{p} = \mathbf{A}^T\boldsymbol{\lambda} - \mathbf{g}, \quad (7)$$

where the residual  $\mathbf{r} = \mathbf{A}\mathbf{x}_0 - \mathbf{b}$  and  $\mathbf{g} = \mathbf{Q}\mathbf{x}_0 + \mathbf{d}$  (the gradient of  $q(\mathbf{x}_0)$ ). Explain how the two equations can be used to solve for  $\boldsymbol{\lambda}$  and  $\mathbf{p}$ .

(See the following page for part (b) of this question.)

(b) Consider the **inequality**-constrained QP:

$$\min_{x \in \mathbb{R}^2} q(x) \text{ subject to } Ax \geq \mathbf{b}, \quad (8)$$

with  $q(x) = \frac{1}{2}x^T Qx + d^T x$  where  $Q = \begin{bmatrix} 2 & 0 \\ 0 & 2 \end{bmatrix}$  and  $d = \begin{bmatrix} -2 \\ -5 \end{bmatrix}$  with

$$A = \begin{bmatrix} 1 & -2 \\ -1 & -2 \\ -1 & 2 \\ 1 & 2 \end{bmatrix} \text{ and } \mathbf{b} = \begin{bmatrix} -2 \\ -6 \\ -2 \\ 2 \end{bmatrix}.$$

- (i) Sketch the diamond-shaped feasible region, labelling the constraints 1–4 in order of their appearance in  $A$  and  $\mathbf{b}$ , e.g. constraint #1 is the region **below** the line  $x - 2y = -2$  and constraint #3 is the region **above** the line  $-x + 2y = -2$ . 2
- (ii) Starting at the vertex  $x_0 = (2, 0)$  where constraints #3 and #4 are binding/active, solve the problem by solving a succession of equality-constrained QPs

$$\min_{x \in \mathbb{R}^2} q(x) \text{ subject to } \hat{A}x = \hat{\mathbf{b}}, \quad (9)$$

where  $\hat{A}$  consists of the rows of  $A$  corresponding to the indices in the working set  $\mathcal{W}$  and similarly for  $\hat{\mathbf{b}}$ . Note that  $\mathbf{r}$  is the residual  $\hat{A}x - \hat{\mathbf{b}}$ .

Use the following steps:

- A. Explain why  $\mathbf{r}$  is the zero vector at the starting point  $x_0$  and calculate  $\mathbf{g}(x_0)$ . (You should find that  $\mathbf{g} = \begin{bmatrix} 2 \\ -5 \end{bmatrix}$ .)

Take the working set  $\mathcal{W} = \{3, 4\}$ , use Eq. 6 in Part (a) of this question to show  $\lambda = \begin{bmatrix} -9/4 \\ -1/4 \end{bmatrix}$  (take  $(\hat{A}Q^{-1}\hat{A}^T)^{-1} = \frac{1}{8} \begin{bmatrix} 5 & -3 \\ -3 & 5 \end{bmatrix}$ ).

Give a simple argument as to why the step  $\mathbf{p}$  **must** be the zero vector. 4

- B. Remove the index corresponding to the most negative multiplier from  $\mathcal{W}$ , update  $\hat{A}$  and  $\hat{\mathbf{b}}$  and resolve for  $\lambda$  and  $\mathbf{p}$  using Eqs. 6 and 7 from Part (a) of this question.

(You should find  $\lambda = -8/5$  and  $\mathbf{p} = \begin{bmatrix} -9/5 \\ 9/10 \end{bmatrix}$ .) 3

C. Determine the steplength  $\alpha$  using the formula

$$\alpha \equiv \min\left(1, \min_{i \in \mathcal{W}^c, \mathbf{a}_i^T \mathbf{p} < 0} \frac{\mathbf{b}_i - \mathbf{a}_i^T \mathbf{x}}{\mathbf{a}_i^T \mathbf{p}}\right). \quad (10)$$

You should find that  $\alpha = \min\{1, 10/9\} = 1$  so no constraint is added to  $\mathcal{W}$ .

3

D. Update  $\mathbf{x}$  to  $\mathbf{x} + \alpha \mathbf{p}$  and recalculate  $\mathbf{r}$  and  $\mathbf{g}$  at the new point  $\mathbf{x} = \begin{bmatrix} 1/5 \\ 9/10 \end{bmatrix}$ . (You should find that  $\mathbf{r}$  is still zero and that  $\mathbf{g} = \begin{bmatrix} -8/5 \\ -16/5 \end{bmatrix}$ .)

3

E. At the next step (**do not perform the calculation**)  $\lambda$  and  $\mathbf{p}$  are found to be  $\lambda = -8/5$  and  $\mathbf{p} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$ . Explain why you should drop the (only) constraint from the working set. Solve the **unconstrained** problem  $\min \frac{1}{2} \mathbf{p}^T \mathbf{Q} \mathbf{p} + \mathbf{g}^T \mathbf{p}$ . You should find  $\mathbf{p} = \begin{bmatrix} 4/5 \\ 8/5 \end{bmatrix}$ .

2

F. At a later iteration,  $\mathcal{W} = \{1\}$  and  $\mathbf{x} = \begin{bmatrix} 7/5 \\ 17/10 \end{bmatrix}$ ,  $\lambda = 4/5$  and  $\mathbf{p} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$ . Explain carefully why the first-order KKT conditions for the inequality-constrained problem Eq. 8 are now satisfied.

2

G. Why is it obvious that the second-order conditions are also satisfied?

1

## Appendix of Results

A The Wolfe conditions for the step length  $\alpha$  in a line search require that

$$f(\mathbf{x}_k + \alpha \mathbf{p}_k) \leq f(\mathbf{x}_k) + c_1 \alpha \mathbf{p}_k^T \mathbf{g}(\mathbf{x}_k), \quad (10a)$$

$$\mathbf{p}_k^T \mathbf{g}(\mathbf{x}_k + \alpha \mathbf{p}_k) \geq c_2 \mathbf{p}_k^T \mathbf{g}(\mathbf{x}_k), \quad (10b)$$

where  $\mathbf{g}(\mathbf{x}) \equiv \nabla f(\mathbf{x})$  and  $0 < c_1 < c_2 < 1$ . The **strong** Wolfe conditions replace (10b) by

$$|\mathbf{p}_k^T \mathbf{g}(\mathbf{x}_k + \alpha \mathbf{p}_k)| \leq c_2 |\mathbf{p}_k^T \mathbf{g}(\mathbf{x}_k)|. \quad (11)$$

B In terms of a “line” function  $\phi(\alpha) \equiv f(\mathbf{x} + \alpha \mathbf{p})$ , the Wolfe conditions for the step length  $\alpha$  in a line search require that

$$\phi(\alpha) \leq \phi(0) + c_1 \alpha \phi'(0), \quad (11a)$$

$$\phi'(\alpha) \geq c_2 \phi'(0), \quad (11b)$$

where  $0 < c_1 < c_2 < 1$ . The **strong** Wolfe conditions replace (11b) by

$$|\phi'(\alpha)| \leq c_2 |\phi'(0)|. \quad (12)$$

C **Theorem 1 (Zoutendijk)** Consider any iteration of the form  $\mathbf{x}_{k+1} = \mathbf{x}_k + \alpha_k \mathbf{p}_k$ , where  $\mathbf{p}_k$  is a descent direction and  $\alpha_k$  satisfies the Wolfe conditions Eqs. 10a and 10b in Appendix A above. Suppose that  $f$  is bounded below in  $\mathbb{R}^n$  and that  $f$  is  $C^1$  in an open set  $\mathcal{N}$  containing the level set  $\mathcal{L} \equiv \{\mathbf{x} : f(\mathbf{x}) \leq f(\mathbf{x}_0)\}$ , where  $\mathbf{x}_0$  is the starting point. Also assume that  $\mathbf{g}(\mathbf{x})$ , the gradient of  $f$ , is Lipschitz continuous on  $\mathcal{N}$ , i.e. there exists a constant  $L$  such that

$$\|\mathbf{g}(\mathbf{x}) - \mathbf{g}(\bar{\mathbf{x}})\| \leq L \|\mathbf{x} - \bar{\mathbf{x}}\|, \quad \text{for all } \mathbf{x}, \bar{\mathbf{x}} \in \mathcal{N}. \quad (13)$$

Then

$$\sum_{k \geq 0} \cos^2 \theta_k \|\mathbf{g}(\mathbf{x}_k)\|^2 < \infty, \quad (14)$$

where  $\theta_k$  is the angle between  $\mathbf{p}_k$  and the steepest descent direction  $-\mathbf{g}(\mathbf{x}_k)$ .

D A sequence  $\{\mathbf{x}_k\}_{k=1}^{\infty}$  that converges to a point  $\mathbf{x}^*$  has quadratic convergence if for some positive constant  $M$

$$\frac{\|\mathbf{x}_{k+1} - \mathbf{x}^*\|}{\|\mathbf{x}_k - \mathbf{x}^*\|^2} \leq M, \quad \text{for all } k \text{ sufficiently large.} \quad (15)$$



## E Algorithm 1 (FR-CGM)

```

begin
  Given  $x_0$ .
  set  $g_0 \leftarrow \nabla f_0, p_0 \leftarrow -g_0, k \leftarrow 0$ ;
  while  $g_k \neq 0$  do
     $\alpha_k \leftarrow$  Result of line search along  $p_k$ ;
     $x_{k+1} \leftarrow x_k + \alpha_k p_k$ ;
     $g_{k+1} \leftarrow \nabla f_{k+1}$ ;
     $\beta_{k+1}^{\text{FR}} \leftarrow \frac{\|g_{k+1}\|^2}{\|g_k\|^2}$ ;
     $p_{k+1} \leftarrow -g_{k+1} + \beta_{k+1}^{\text{FR}} p_k$ ;
     $k \leftarrow k + 1$ ;
  end (while)
end

```

F The Sherman-Morrison-Woodbury (SMW) formula states that if a square non-singular matrix  $A$  is updated by

$$A' = A + \Delta A, \quad \text{where } \Delta A = RST^T$$

where  $R, T$  are  $n \times p$  matrices for  $1 < p < n$  and  $S$  is  $p \times p$  then

$$A'^{-1} \equiv (A + \Delta A)^{-1} = A^{-1} - A^{-1}R U^{-1} T^T A^{-1}, \quad (16)$$

where  $U = S^{-1} + T^T A^{-1} R$ .

G Let  $x^*$  be a local minimum of the equality-constrained optimisation problem defined in App. H . For each positive integer  $k$ , define a **penalty function**

$$F^k(x) = f(x) + \frac{k}{2} \|c(x)\|^2 + \frac{\alpha}{2} \|x - x^*\|^2, \quad (17)$$

where  $\alpha > 0$  is arbitrary.

H The general equality constrained optimisation problem is:

$$\min_{x \in \mathbb{R}^n} f(x) \quad \text{subject to} \quad c_i(x) = 0, i \in \mathcal{E}. \quad (18)$$

I The general inequality constrained optimisation problem is:

$$\min_{x \in \mathbb{R}^n} f(x) \quad \text{subject to} \quad \begin{cases} c_i(x) = 0 & i \in \mathcal{E}, \\ c_i(x) \geq 0 & i \in \mathcal{I}. \end{cases} \quad (19)$$

J The first-order KKT necessary conditions for a point  $\mathbf{x}^*$  with optimal multipliers  $\lambda^*$  to be a local solution of an inequality-constrained minimisation problem are as follows: Suppose that  $\mathbf{x}^*$  is a local solution of a general constrained optimisation (as in App. I) problem and that the LICQ holds at  $\mathbf{x}^*$  (i.e. the active constraint gradients are linearly independent). Then there is a Lagrange multiplier vector  $\lambda^*$ , with components  $\lambda_i^*$ ,  $i \in \mathcal{E} \cup \mathcal{I}$ , such that the following conditions are satisfied at  $(\mathbf{x}^*, \lambda^*)$ :

$$\nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x}^*, \lambda^*) = 0, \quad (20a)$$

$$\mathbf{c}_i(\mathbf{x}^*) = 0, \quad \text{for all } i \in \mathcal{E}, \quad (20b)$$

$$\mathbf{c}_i(\mathbf{x}^*) \geq 0, \quad \text{for all } i \in \mathcal{I}, \quad (20c)$$

$$\lambda_i^* \geq 0, \quad \text{for all } i \in \mathcal{I}, \quad (20d)$$

$$\lambda_i^* \mathbf{c}_i(\mathbf{x}^*) = 0, \quad \text{for all } i \in \mathcal{E} \cup \mathcal{I} \quad (20e)$$

where  $\mathcal{L}(\mathbf{x}, \lambda) \equiv f(\mathbf{x}) - \sum_{i \in \mathcal{E} \cup \mathcal{I}} \lambda_i \mathbf{c}_i(\mathbf{x})$ .

K The (revised) Simplex algorithm may be stated as:

### Algorithm 2

Given  $\mathcal{B}, \mathcal{N}, \mathbf{x}_B = \mathbf{B}^{-1} \mathbf{b} \geq 0, \mathbf{x}_N = 0$ ;

Solve  $\mathbf{B}^T \lambda = \mathbf{c}_B$  for  $\lambda$

$\mathbf{s}_N = \mathbf{c}_N - \mathbf{N}^T \lambda$ ; ( pricing )

if  $\mathbf{s}_N \geq 0$

    then STOP( optimal point found )

fi

Select  $q \in \mathcal{N}$  with  $s_q < 0$  as the entering index

Solve  $\mathbf{B} \mathbf{d} = \mathbf{A}_q$  for  $\mathbf{d}$ ;

if  $\mathbf{d} \leq 0$

    then STOP( problem is unbounded )

fi

Calculate  $\mathbf{x}_q^+ = \min_{i | d_i > 0} (\mathbf{x}_B)_i / d_i$  and use  $\mathbf{p}$  to denote the minimizing  $i$ ;

Update  $\mathbf{x}_B^+ = \mathbf{x}_B - \mathbf{d} \mathbf{x}_q^+, \mathbf{x}_N^+ = (0, \dots, 0, \mathbf{x}_q^+, 0, \dots, 0)^T$ ;

Change  $\mathcal{B}$  by adding  $q$  and removing the basic variable corresponding to column  $\mathbf{p}$  of  $\mathcal{B}$ .