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College of Informatics and Electronics

**END OF SEMESTER ASSESSMENT PAPER**

MODULE CODE: MS4327

SEMESTER: Spring 2005

MODULE TITLE: Optimisation

DURATION OF EXAMINATION: 2 1/2 hours

LECTURER: Dr. J. Kinsella

PERCENTAGE OF TOTAL MARKS: 70%

EXTERNAL EXAMINER: Prof. J. King

**INSTRUCTIONS TO CANDIDATES: Answer four questions correctly for full marks; 70%.  
See the Appendix at the end of the paper for some useful results..**

- 1 (a) Prove Zoutendijk's Theorem (Theorem 2 in Appendix B) 15
- (b) Suppose that the search directions  $p_k$  are generated using a Newton-like method:  $p_k = -B_k^{-1}g(x_k)$  where  $B_k$  is symmetric and positive definite. Show that if  $\|B_k\|\|B_k^{-1}\| \leq M$  for all  $k$  then  $\cos \theta_k \geq 1/M$  where  $\theta_k$  is as defined above. 8
- (c) Use Zoutendijk's Theorem above to show that in this case  $\lim_{k \rightarrow \infty} \|g(x_k)\| = 0$ . 2
- 2 The Dogleg Trust Region method is described in Appendix C.
- (a) Show that if  $B$  is positive definite then
- (i)  $\|\tilde{p}(\tau)\|$  is an increasing function of  $\tau$  and 5
- (ii)  $m(\tilde{p}(\tau))$  is a decreasing function of  $\tau$ . 5
- (b) Give a detailed argument which demonstrates that the path  $\tilde{p}(\tau)$  intersects the trust region boundary  $\|p\| = \Delta$  at exactly one point if  $\|\tilde{p}(2)\| \equiv \|p_B\| \geq \Delta$  and nowhere otherwise. 2
- (c) What value should be selected for  $p$  if  $\|p_B\| \leq \Delta$ ? 3
- (d) In the case where the vector  $p$  is chosen to be on the boundary, explain carefully how the appropriate value of the parameter  $\tau$  is chosen. 5
- (e) Finally, if  $p_U = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$ ,  $p_B = \begin{bmatrix} 1.5 \\ 0.5 \end{bmatrix}$  and  $\Delta = 1.5$ , find the appropriate value of  $\tau$ . 5
- 3 Given the Preliminary form of the Linear Conjugate Gradient Method (Appendix D); prove 25

**Theorem 1** Suppose that the  $k^{\text{th}}$  iterate generated by the conjugate gradient method Alg. 1 is not the solution point  $x^*$ . The following properties hold:

$$p_k^T A p_i = 0, \quad \text{for } i = 0, \dots, k-1. \quad (1)$$

$$r_k^T r_i = 0, \quad \text{for } i = 0, \dots, k-1, \quad (2)$$

4 Given the Dai-Yuan version (Appendix F) of the Non-linear Conjugate Gradient Method (Appendix E) prove that the method has the property that under the same assumptions as those of Zoutendijk's Theorem (Appendix B), in particular

- the weak Wolfe conditions hold (Appendix A)
- $f$  is bounded below
- with the exception that we need not assume that the  $p_k$  are descent directions

we have that the algorithm either stops at a stationary point ( $\|g_k\| = 0$ ) or  $\liminf \|g_k\| = 0$ .

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5 The BFGS method updates an estimate  $H_k$  of the Inverse Hessian in the form:

$$H_{k+1} = (I - \gamma_k s_k y_k^T) H_k (I - \gamma_k y_k s_k^T) + \gamma_k s_k s_k^T. \quad (3)$$

Here  $s_k = x_{k+1} - x_k$ ,  $y_k = \nabla f(x_{k+1}) - \nabla f(x_k)$  and  $\gamma_k = s_k^T y_k$ . Derive this update rule by solving the problem:

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$$\min_H \|H - H_k\| \quad (3a)$$

$$\text{subject to } H = H^T, H y_k = s_k. \quad (3b)$$

Use the "weighted Frobenius norm":  $\|A\|_W \equiv \|W^{\frac{1}{2}} A W^{\frac{1}{2}}\|_F$ , where  $\|C\|_F^2 \equiv \sum_{i=1}^n \sum_{j=1}^n C_{ij}^2$ . The weight  $W$  is an unspecified positive definite matrix satisfying  $W s_k = y_k$ .

6 Consider the equality-constrained Quadratic Program (Q.P.):

$$\min_x q(x) = \frac{1}{2}x^T Gx + x^T d, \quad (4a)$$

$$\text{subject to } a_i^T x = b, \quad i = 1, \dots, k \quad (4b)$$

- (a) Writing the set of  $k$  equality constraints (4b) as the matrix equation  $Ax - b = 0$ , show that if  $x^*$  is a local minimum then the KKT conditions (Appendix G) require that there must be a vector  $\lambda^*$  of Lagrange multipliers such that the following system of equations is satisfied: 5

$$\begin{bmatrix} G & -A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} x^* \\ \lambda^* \end{bmatrix} = \begin{bmatrix} -d \\ b \end{bmatrix} \quad (5)$$

- (b) If we write  $x^* = x + p$ , where  $x$  is an estimate of the solution and  $p$  the required step to the solution, show that (5) can be re-written as: 2

$$\begin{bmatrix} G & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} -p \\ \lambda^* \end{bmatrix} = \begin{bmatrix} g \\ c \end{bmatrix} \quad (6)$$

where  $c = Ax - b$ ,  $g = Gx + d$  (the gradient of  $q(x)$ ) and  $p = x^* - x$ .

- (c) Show that this block matrix equation can be reduced to the problem of solving

$$\left( AG^{-1}A^T \right) \lambda^* = \left( AG^{-1}g - c \right). \quad (7)$$

for  $\lambda^*$  and then solving

$$-Gp + A^T \lambda^* = g \quad (8)$$

for  $p$ . 5

- (d) Given  $q(x) = \frac{1}{2}x_1^2 + 2x_1x_2 - x_2^2 + 3x_1 + 5x_2$ . and the equality constraint  $x_1 - 2x_2 = 3$ , solve the QP using the method you derived in (c). 8

- (e) Finally check that your solution satisfies the second-order conditions (Appendix H). 5

## Appendix of Results

A The Wolfe conditions for the step length  $\alpha$  in a line search require that

$$f(x_k + \alpha p_k) \leq f(x_k) + c_1 \alpha p_k^T g(x_k), \quad (8a)$$

$$p_k^T g(x_k + \alpha p_k) \geq c_2 p_k^T g(x_k) \quad (8b)$$

where  $g(x) \equiv \nabla f(x)$  and  $0 < c_1 < c_2 < 1$ . The strong Wolfe conditions replace (8b) by

$$|p_k^T g(x_k + \alpha p_k)| \leq c_2 |p_k^T g(x_k)|. \quad (9)$$

**B Theorem 2 (Zoutendijk)** Consider any iteration of the form  $x_{k+1} = x_k + \alpha_k p_k$ , where  $p_k$  is a descent direction and  $\alpha_k$  satisfies the Wolfe conditions Eqs. 8a and 8b in Appendix A above. Suppose that  $f$  is bounded below in  $\mathbb{R}^n$  and that  $f$  is  $C^1$  in an open set  $\mathcal{N}$  containing the level set  $\mathcal{L} \equiv \{x : f(x) \leq f(x_0)\}$ , where  $x_0$  is the starting point. Also assume that  $g(x)$ , the gradient of  $f$ , is Lipschitz continuous on  $\mathcal{N}$ , i.e. there exists a constant  $L$  such that

$$\|g(x) - g(\bar{x})\| \leq L \|x - \bar{x}\|, \quad \text{for all } x, \bar{x} \in \mathcal{N}. \quad (10)$$

Then

$$\sum_{k \geq 0} \cos^2 \theta_k \|g(x_k)\|^2 < \infty, \quad (11)$$

where  $\theta_k$  is the angle between  $p_k$  and the steepest descent direction  $-g(x_k)$ .

C The Trust Region Method is based on the problem:

$$\min_{p \in \mathbb{R}^n} m(p) \equiv f_0 + g^T p + \frac{1}{2} p^T B p, \quad \text{such that } \|p\| \leq \Delta, \quad (12)$$

where  $f_0$  is a fixed scalar,  $g$  a fixed vector in  $\mathbb{R}^n$ ,  $B$  a fixed  $n \times n$  matrix and  $\Delta$  a fixed positive scalar. The “dogleg” method finds an approximate solution to (12) by replacing the (unknown) curved trajectory for  $p^*(\Delta)$  with a path consisting of two line segments. The first line segment runs from the starting point to the unconstrained minimiser along the steepest descent direction defined by

$$p^U = -\frac{g^T g}{g^T B g} g \quad (13)$$

while the second line segment runs from  $p^U$  to  $p^B \equiv -B^{-1}g$ . We can define the trajectory as a path  $\tilde{p}(\tau)$  parameterised by  $\tau$  as follows:

$$\tilde{p}(\tau) = \begin{cases} \tau p^U, & 0 \leq \tau \leq 1, \\ p^U + (\tau - 1)(p^B - p^U), & 1 \leq \tau \leq 2. \end{cases} \quad (14)$$

## D Algorithm 1 (CGM—Preliminary Version)

**begin**Given  $x_0$ .**set**  $r_0 \leftarrow Ax_0 - b, p_0 \leftarrow -r_0, k \leftarrow 0$ ;**while**  $r_k \neq 0$  **do** $\alpha_k \leftarrow -\frac{r_k^T p_k}{p_k^T A p_k}$ ; **minimises residual along**  $p_k$  $x_{k+1} \leftarrow x_k + \alpha_k p_k$ ; $r_{k+1} \leftarrow Ax_{k+1} - b$ ; $\beta_{k+1} \leftarrow \frac{r_{k+1}^T A p_k}{p_k^T A p_k}$ ; $p_{k+1} \leftarrow -r_{k+1} + \beta_{k+1} p_k$ ; **so**  $p_{k+1}$  **conjugate to**  $p_k$  $k \leftarrow k + 1$ ;**end** (while)**end**

E The original Fletcher-Reeves Conjugate Gradient Method is given in the following algorithm:

**Algorithm 2****begin**Given  $x_0$ .**set**  $r_0 \leftarrow \nabla f_0, p_0 \leftarrow -r_0, k \leftarrow 0$ ;**while**  $r_k \neq 0$  **do** $\alpha_k \leftarrow$  **Result of line search along**  $p_k$ ; $x_{k+1} \leftarrow x_k + \alpha_k p_k$ ; $r_{k+1} \leftarrow \nabla f_{k+1}$  $\beta_{k+1}^{FR} \leftarrow \frac{r_{k+1}^T r_{k+1}}{r_k^T r_k} \equiv \frac{\|r_{k+1}\|^2}{\|r_k\|^2}$ ; $p_{k+1} \leftarrow -r_{k+1} + \beta_{k+1}^{FR} p_k$ ; $k \leftarrow k + 1$ ;**end** (while)**end**

F The Dai-Yuan (DY) version of the Conjugate Gradient Method replaces  $\beta^{FR}$  in Alg. 2 by

$$\beta_{k+1}^{DY} = \frac{\|g_{k+1}\|^2}{p_k^T y_k}, \quad (15)$$

where  $y_k \equiv g_{k+1} - g_k$

G The first-order KKT necessary conditions for a point  $x^*$  with optimal multipliers  $\lambda^*$  to be a local solution of an equality-constrained minimisation problem

$$\min f(x) \text{ subject to } c_i(x) = 0, i = 1, \dots, k$$

are

$$\nabla_x \mathcal{L}(x^*, \lambda^*) = 0, \quad (16a)$$

$$c_i(x^*) = 0, \quad i = 1, \dots, k, \quad (16b)$$

where  $\mathcal{L}(x, \lambda) = f(x) - \sum_{i=1}^k \lambda_i c_i(x)$

H The second-order KKT necessary conditions for a point  $x^*$  with optimal multipliers  $\lambda^*$  to be a local solution of the above equality-constrained minimisation problem are that  $d^T \nabla^2 f(x^*) d \geq 0$  for all stationary directions  $d$  (directions that satisfy  $d^T \nabla c_i(x^*) = 0$ ).