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OLLSCOIL LUIMNIGH

Faculty of Science & Engineering
Department of Mathematics & Statistics

END OF SEMESTER ASSESSMENT PAPER

MODULE CODE: MS4105

SEMESTER: Autumn 2008/9

MODULE TITLE: Linear Algebra 2

DURATION OF EXAMINATION: 2 1/2 hours

LECTURER: Dr. J. Kinsella

PERCENTAGE OF TOTAL MARKS: 100%

EXTERNAL EXAMINER: Prof. J. Flavin

INSTRUCTIONS TO CANDIDATES: Answer four questions correctly for full marks, 100%.

- 1 (a) Let V be a finite-dimensional vector space and let $\{v_1, \dots, v_n\}$ be a basis for V . Prove **one** of the following propositions (**A or B**): 13%
- (A) If a set of vectors in V has more than n vectors, it is linearly dependent.
- (B) If a set of vectors in V has fewer than n vectors, it cannot span V . You may assume that a homogeneous linear system with more unknowns than equations has infinitely many solutions.
- (b) Explain briefly how results (A) and (B) lead to a definition for the dimension of a vector space. 2%
- (c) If u and v are vectors in a real inner product space then prove the Cauchy-Schwarz inequality: 6%

$$|\langle u, v \rangle| \leq \|u\| \|v\|. \quad (1)$$

- (d) Use the Cauchy-Schwarz inequality to derive the Triangle Inequality: 4%

$$\|u + v\| \leq \|u\| + \|v\| \quad (2)$$

- 2 (a) Suppose that an $n \times n$ real matrix A has full rank (the columns of A are linearly independent). Show that A is invertible. (You may assume that the columns of A span \mathbb{R}^n .) 6%
- (b) Suppose that the complex matrix S is skew-hermitian ($S^* = -S$ where S^* is the hermitian conjugate of S). Show that the matrix

$$Q = (I - S)^{-1}(I + S)$$

is unitary ($Q^*Q = I$). 5%

- (c) Does a unitary matrix always have the property $QQ^* = I$? Explain briefly. 1%
- (d) Show that for any vector $x \in \mathbb{C}^m$; 2%+2%
- (i) $\|x\|_\infty \leq \|x\|_2$,
- (ii) $\|x\|_2 \leq \sqrt{m}\|x\|_\infty$,
- (e) Using the general definition of an induced matrix norm

$$\|A\| = \sup_{x \in \mathbb{C}^n, x \neq 0} \frac{\|Ax\|}{\|x\|},$$

show using the results from 2.(d) that for any $m \times n$ complex matrix A ; 5%+4%

- (i) $\|A\|_\infty \leq \sqrt{n}\|A\|_2$,
- (ii) $\|A\|_2 \leq \sqrt{m}\|A\|_\infty$.

- 3 (a) Show that for every $m \times n$ complex matrix A we can write: 10%

$$A_1 \equiv U^*AV = \begin{bmatrix} \sigma & 0 \\ 0 & B \end{bmatrix} \quad (3)$$

where $\sigma = \|A\|_2$, B is an $(m-1) \times (n-1)$ matrix, $U = [y_0 \ U_1]$ and $V = [x_0 \ V_1]$, the unit vector x_0 satisfies $\|Ax_0\| = \|A\|$ and finally $y_0 = \frac{Ax_0}{\|A\|}$. The matrices U_1 and V_1 are chosen so that U and V are unitary $m \times m$ and $n \times n$ respectively.

- (b) Using part (a) prove by induction on the size of A that every $m \times n$ complex matrix A has a Singular Value Decomposition (SVD) $A = U\Sigma V^*$ where U is $m \times m$ unitary, V is $n \times n$ unitary and Σ , is an $m \times n$ diagonal matrix of singular values. 8%
- (c) Use the SVD to show that the row and column ranks of any $m \times n$ matrix are equal to r , the number of non-zero singular values. (Assume that $m > n$.) 7%

- 4 (a) For any vector $v \in \mathbb{C}^p$, let the matrix $H = I - 2P_v$ (where $P = \frac{vv^*}{v^*v}$);

- (i) Show with a sketch that the effect on an arbitrary vector $x \in \mathbb{C}^p$ of left-multiplying x by H is to reflect x in the normal to v in the x - v plane. 3%
- (ii) Find the choices of vector v that result in the **Householder reflection** of x , Hx , being a multiple of e_1 where $e_1 \in \mathbb{C}^p$ is a vector of zeroes with one in the first position. 7%
- (iii) Which of the two choices found should be used and why? 1%
- (iv) Given the vector $x = (-1, 3, 1, -2)^T \in \mathbb{R}^4$, find the appropriate choice of v that ensures that Hx is a multiple of e_1 . 2%
- (v) With this choice of v , what is the value of Hx ?

- (b) Explain the operation of the following algorithm (Alg. 0.1) which takes as its input an arbitrary matrix $m \times n$ complex matrix A .
- (i) Explain how the choice of v_k in lines 3 & 4 relates to that found in part (a). 1%
 - (ii) Explain the effect of line 5 and relate it to the Householder reflection in part (a). 3%
 - (iii) What special structure will the matrix A have after the algorithm has completed? 1%

Algorithm 0.1 *Householder QR Factorisation*

```

1   for  $k = 1$  to  $n$ 
2        $x = A_{k:m,k}$ 
3        $v_k = x + \text{sign}(x_1)\|x\|e_1$ 
4        $v_k = v_k/\|v_k\|$ 
5        $A_{k:m,k:n} = A_{k:m,k:n} - 2v_k(v_k^*A_{k:m,k:n})$ 
6   end

```

- (c) The work done in Alg. 0.1 is dominated by the (implicit) inner loop $j=k:n$ over the columns of the submatrix $A_{k:m,k:n}$ in line 5. Show that the operation count for the algorithm is $\approx 2mn^2 - 2/3n^3$. 7%

5 For any $m \times m$ matrix A , Gauss Elimination without pivoting consists of:

Algorithm 0.2 *Gauss Elimination Without Pivoting — in words*

```

1   for  $k = 1$  to  $m - 1$ 
2       Add suitable multiples of row  $k$  to the rows beneath
3       to introduce zeroes below the main diagonal in column  $k$ .
4   end

```

- (a) Show that each iteration of the above algorithm can be effected by left-multiplying A by a matrix $L_k = I - \ell_k e_k^*$ where ℓ_k is the vector of **multipliers** for the k^{th} column (the first k entries of ℓ_k are 0) and e_k is a vector in \mathbb{C}^n with one in the k^{th} position and zeroes elsewhere. Give a simple formula for the non-zero entries of ℓ_k . 6%
- (b) Show that each L_k can be inverted by negating its sub-diagonal elements. 3%
- (c) Show that the product of the $m - 1$ matrices L_k^{-1} in increasing order of k is the identity matrix with the non-zero sub-diagonal elements of the L_k^{-1} inserted in the appropriate places. 3%
- (d) What special structure does A have after the algorithm has completed? 1%

- (e) Explain the operation of the following implementation of Gauss Elimination, algorithm (Alg. 0.3) which takes as its input an arbitrary matrix $m \times m$ complex matrix A , in particular, how does the inner loop (lines 3–6) correspond to multiplication by L_k ? 3%

Algorithm 0.3 *Gauss Elimination without Pivoting*

```

1      U = A, L = I
2      for k = 1 to m - 1
3          for j = k + 1 to m
4               $l_{jk} = \frac{u_{jk}}{u_{kk}}$ 
5               $u_{j,k:m} = u_{j,k:m} - l_{jk}u_{k,k:m}$ 
6          end
7      end

```

- (f) The work done in Alg. 0.3 is dominated by Line 5. Show that the operation count for the algorithm is $\approx 2/3m^3$. 6%
- (g) Explain briefly how partial pivoting works, when it must be used and why it is always a good strategy. 3%
- 6 (a) Show that every $m \times m$ matrix A has the property that: 4%

$$Q_0^* A Q_0 = \begin{bmatrix} \lambda_1 & v_1^* \\ 0 & A_1 \end{bmatrix}, \quad (4)$$

where $Q_0 = [q_1 \ q'_2 \ \dots \ q'_n]$, q_1 is an eigenvector of A , λ_1 is the corresponding eigenvalue and the vectors q'_2, \dots, q'_n are chosen so that $\{q_1, q'_2, \dots, q'_n\}$ form an orthonormal basis for \mathbb{C}^n . (Also $v_1 \in \mathbb{C}^{n-1}$ and $A_1 \in \mathbb{C}^{(m-1) \times (m-1)}$.)

- (b) Use induction to prove that every $m \times m$ complex matrix A has a Schur decomposition: 6%

$$A = QTQ^* \quad (5)$$

where Q is unitary and T is upper triangular.

- (c) Show that, given the Schur decomposition $A = QTQ^*$, the eigenvalues of A appear on the main diagonal of T . 8%
- (d) Explain briefly why an **exact** method (one that calculates the eigenvalues in a finite number of steps in exact arithmetic) for calculating the eigenvalues of a matrix A is impossible. 1%

- (e) The following algorithm (Alg. 0.4) reduces any $m \times m$ matrix A to Hessenberg form using Householder reflections by multiplying A successively on the left by Q_k^* and on the right by Q_k , $k = 1, \dots, m - 2$.

The matrix $Q_k = \begin{bmatrix} I_k & 0 \\ 0 & H_k \end{bmatrix}$, I_k is the $k \times k$ identity matrix and H_k is the Householder reflector (see Q.4(a)) that transforms the vector $x_k = A_{k+1:m,k}$ into a multiple of e_1 at each iteration.

Explain carefully how lines 5 and 6 of Alg. 0.4 correspond to left-multiplication by Q_k^* and right-multiplication by Q_k respectively (note that the matrices Q_k are hermitian, i.e. $Q_k^* = Q_k$):

6%

Algorithm 0.4 *Householder Reduction to Hessenberg Form*

```

1      for  $k = 1$  to  $m - 2$ 
2           $x = A_{k+1:m,k}$ 
3           $v_k = x + \text{sign}(x_1) \|x\| e_1$ 
4           $v_k = v_k / \|v_k\|$ 
5           $A_{k+1:m,k:m} = A_{k+1:m,k:m} - 2v_k (v_k^* A_{k+1:m,k:m})$ 
6           $A_{1:m,k+1:m} = A_{1:m,k+1:m} - 2(A_{1:m,k+1:m} v_k) v_k^*$ 
7      end

```